

Fourth and Final Examination *solutions*

Math 320, David Pierce

June 11 (Saturday), 2011, at 13:30 in M-04

Instructions. Each of the five numbered problems is worth 8 points. As usual, write solutions on separate sheets, keeping *this* sheet. Then enjoy the holiday!

Problem 1. Write down formulas defining the following classes. (Use only the symbols \in , \neg , $($, \Rightarrow , $)$, and \exists ; variables; and the constant a .)

a) $\mathcal{P}(a)$

b) $\bigcup a$

Solution.

a) $x \subseteq a$, that is, $\forall y (y \in x \Rightarrow y \in a)$, that is,

$$\neg \exists y \neg (y \in x \Rightarrow y \in a).$$

b) $\exists y (y \in a \wedge x \in y)$.

Remark. The formulas *defining* the classes are as given. Then for example the class $\mathcal{P}(a)$ itself is

$$\{x: \neg\exists y \neg(y \in x \Rightarrow y \in a)\}.$$

Problem 2. *Prove or disprove:*

- a) *Every set is a class.*
- b) *Every class is a set.*

Solution.

- a) Every set a is the class $\{x: x \in a\}$.
- b) Not every class is a set. Indeed, the class $\{x: x \notin x\}$ is not a set, for if it were the set a , then

$$\forall x (x \in a \Leftrightarrow x \notin x),$$

and in particular

$$a \in a \Leftrightarrow a \notin a,$$

which is a contradiction.

Problem 3. *Perform the following ordinal computations, giving the answers in Cantor normal form.*

Solution.

- a) $3 \cdot (\omega + 4) = 3 \cdot \omega + 3 \cdot 4 = \omega + 12$
- b) $(\omega + 4) \cdot 3 = \omega \cdot 3 + 4$
- c) $(\omega + 5)^2 = (\omega + 5) \cdot (\omega + 5) = (\omega + 5) \cdot \omega + (\omega + 5) \cdot 5 = \omega^2 + \omega \cdot 5 + 5$

$$\text{d) } 9^{\omega+2} = 9^\omega \cdot 9^2 = \omega \cdot 81$$

$$\text{e) } (\omega+5)^{\omega+2} = (\omega+5)^\omega \cdot (\omega+5)^2 = \omega^\omega \cdot (\omega^2 + \omega \cdot 5 + 5) = \omega^{\omega+2} + \omega^{\omega+1} \cdot 5 + \omega^\omega \cdot 5$$

$$\text{f) } (\omega^\omega)^{\omega^\omega} = \omega^{\omega \cdot \omega^\omega} = \omega^{\omega^{1+\omega}} = \omega^{\omega^\omega}$$

$$\text{g) } (\omega^{\omega^\omega})^{\omega^\omega} = \omega^{\omega^\omega \cdot \omega^\omega} = \omega^{\omega^{\omega \cdot 2}}$$

$$\text{h) } 6^{\omega^{1330}} = (6^\omega)^{\omega^{1329}} = \omega^{\omega^{1329}}$$

Problem 4. Prove, for all ordinals α , β , and γ such that $\alpha > 1$,

$$\alpha^{\beta+\gamma} = \alpha^\beta \cdot \alpha^\gamma \quad (*)$$

Use the recursive definitions, and normality, of $x \mapsto \alpha^x$, $x \mapsto \beta+x$, and $x \mapsto \delta \cdot x$ (where $\delta > 0$). You may use other known ordinal identities, besides $(*)$ itself.

Solution. We use induction on γ . Since

$$\alpha^{\beta+0} = \alpha^\beta = \alpha^\beta \cdot 1 = \alpha^\beta \cdot \alpha^0,$$

the claim holds when $\gamma = 0$. Suppose the claim holds when $\gamma = \delta$. Then

$$\alpha^{\beta+\delta'} = \alpha^{(\beta+\delta)'} = \alpha^{\beta+\delta} \cdot \alpha = \alpha^\beta \cdot \alpha^\delta \cdot \alpha = \alpha^\beta \cdot \alpha^{\delta'},$$

so the claim holds when $\gamma = \delta'$.

Suppose finally δ is a limit, and the claim holds when $\gamma < \delta$.

Then

$$\begin{aligned}
 \alpha^{\beta+\delta} &= \alpha^{\sup_{\gamma<\delta}(\beta+\gamma)} && \text{[by definition of } x \mapsto \beta + x\text{]} \\
 &= \sup_{\gamma<\delta} \alpha^{\beta+\gamma} && \text{[by normality of } x \mapsto \alpha^x\text{]} \\
 &= \sup_{\gamma<\delta} (\alpha^\beta \cdot \alpha^\gamma) && \text{[by inductive hypothesis]} \\
 &= \alpha^\beta \cdot \sup_{\gamma<\delta} \alpha^\gamma && \text{[by normality of } x \mapsto \alpha^\beta \cdot x\text{]} \\
 &= \alpha^\beta \cdot \alpha^\delta, && \text{[by definition of } x \mapsto \alpha^x\text{]}
 \end{aligned}$$

so the claim holds when $\gamma = \delta$.

Problem 5. Define the function $\alpha \mapsto V_\alpha$ on **ON** by

$$V_0 = 0, \quad V_{\alpha+1} = \mathcal{P}(V_\alpha), \quad V_\beta = \bigcup_{\alpha<\beta} V_\alpha,$$

where β is a limit. Find $\text{card}(V_\alpha)$ in the following cases. Your answer should be a natural number, an aleph \aleph_β , or a beth \beth_γ .

- | | |
|------------------------------|--------------------------------------|
| a) $\alpha \in \omega$ | e) $\alpha = \omega \cdot 11 + 2011$ |
| b) $\alpha = \omega$ | f) $\alpha = \omega^2$ |
| c) $\alpha = \omega + 320$ | g) $\alpha = \aleph_1$ |
| d) $\alpha = \omega \cdot 6$ | h) $\alpha = \beth_1$ |

Solution.

- a) $\text{card}(V_0) = 0$, and $\text{card}(V_{k+1}) = 2^{\text{card}(V_k)}$ if $k \in \omega$.
b) $\text{card}(V_\omega) = \sup_{k \in \omega} \text{card}(V_k) = \aleph_0$.

- c) $\text{card}(V_{\omega+1}) = 2^{\text{card}(V_\omega)} = 2^{\aleph_0} = 2^{\beth_0} = \beth_1$, and in general

$$\text{card}(V_{\omega+k}) = \beth_k$$

if $k \in \omega$; in particular, $\text{card}(V_{320}) = \beth_{320}$.

- d) $\text{card}(V_{\omega \cdot 2}) = \sup_{k \in \omega} \text{card}(V_{\omega+k}) = \sup_{k \in \omega} \beth_k = \beth_\omega$, and in general

$$\text{card}(V_{\omega \cdot (n+1)}) = \beth_{\omega \cdot n}$$

if $n \in \omega$; in particular, $\text{card}(V_{\omega \cdot 6}) = \beth_{\omega \cdot 5}$.

- e) In general,

$$\text{card}(V_{\omega+\alpha}) = \beth_\alpha \quad (\dagger)$$

for all ordinals α ; in particular, $\text{card}(V_{\omega \cdot 11 + 2011}) = \beth_{\omega \cdot 10 + 2011}$.

- f) Since $\omega^2 = \omega + \omega^2$, we have $\text{card}(V_{\omega^2}) = \beth_{\omega^2}$.

- g) $\text{card}(V_{\aleph_1}) = \beth_{\aleph_1}$

- h) $\text{card}(V_{\beth_1}) = \beth_{\beth_1}$

Remark. The rule (\dagger) can be proved by induction, but this was not required. Note the resemblance to the rule for powers of natural numbers, which can be written as

$$n^{\omega^{1+\alpha}} = n^{\omega \cdot \omega^\alpha} = (n^\omega)^{\omega^\alpha} = \omega^{\omega^\alpha}.$$

where $1 < n < \omega$.

Scores

	EA	PC	AF	Mİ	OŞ	NT	ÖT
1	—	3	7	—	1	7	0
2	—	3	7	0	4	7	0
3	8	7	7	—	6	7	6
4	5	8	7	—	6	8	6
5	0	2	4	—	2	7	0
	13	23	32	0	19	36	12